

Linear Programs and Robust Hedging Problems

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Linear Programs

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Introduction to Linear Programs

- A linear program (LP) is an optimization problem with linear objective function and linear constraints.
- Linear programs can be solved very efficiently (i.e. Simplex method by G.B. Dantzig 1947).
- There are different types of linear programs:
 - ▷ Finite dimensional LPs: finite number of decision variables and constraints;
 - ▷ Semi-infinite LPs: either number of decision variables or constraints is infinite (Polynomial approximation);
 - ▷ Infinite-dimensional LPs: both decision variables and constraints are infinitely dimensional (Optimal transport);
 - ▷ Continuous LPs: linear optimal control problem with linear state constraints (Bottleneck Problem - Bellman 1957).

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Example of a Linear Program in Finite Dimensions

Optimal manufacturing: given a production facility where

- ▶ n is the number of production lines ($i = 1, \dots, n$);
- ▶ m is the number of different products produced on each line ($j = 1, \dots, m$);
- ▶ $x_i \geq 0$ is the level at which each line can be operated;
- ▶ w_j is the revenue collected from producing a unit of product;
- ▶ c_i is the cost of production per line if operated per level;
- ▶ a_{ij} is the yield of each product on each line;
- ▶ b_j is the required output per product.

Example of a Linear Program in Finite Dimensions

- Objective is to produce a given number of products of each category at a minimal cost

$$\min \langle x, c \rangle_{\mathbb{R}^n} \quad \text{s.t.} \quad Ax = b, \quad x \geq 0,$$

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- Can ask a related question of maximizing revenue per unit of production given b units of different products

$$\max \langle b, w \rangle_{\mathbb{R}^m} \quad \text{s.t.} \quad A^* w - c \geq 0, \quad w \in \mathbb{R}^m,$$

where matrix $A^* = A^T$.

Definition of LP

Given two dual pairs of vector spaces (X, Y) and (Z, W) endowed with bilinear forms denoted by $\langle \cdot, \cdot \rangle_{XY}$ and $\langle \cdot, \cdot \rangle_{ZW}$

- Equality constrained problem (EP):

$$\tilde{\mathcal{P}} := \inf \langle x, c \rangle_{XY} \quad \text{s.t.} \quad Ax = b, \quad x \geq 0,$$

where $c \in Y$, $b \in Z$ are given and $A : X \rightarrow Z$ is a linear map.

- Dual equality constrained problem (EP*):

$$\tilde{\mathcal{D}} := \sup \langle b, w \rangle_{ZW} \quad \text{s.t.} \quad A^*w - c \geq 0, \quad w \in W,$$

where $A^* : W \rightarrow Y$ is the adjoint of A such that

$$\langle Ax, w \rangle_{ZW} = \langle x, A^*w \rangle_{XY}.$$

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Weak and strong duality in LPs

- *Feasible solution*: if a decision variable satisfies constraints, it is feasible;
- *Weak duality*: if both the primal and the dual programs have feasible solutions then $\tilde{\mathcal{P}} \leq \tilde{\mathcal{D}}$;
- *Strong duality*: the primal program and its dual have the same value, i.e. $\tilde{\mathcal{P}} = \tilde{\mathcal{D}}$;
- Strong duality always holds for finite dimensional programs;
- Strong duality does not always hold in semi-infinite or infinite dimensional programs – *Duality gap*.

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Interior point conditions for absence of duality gaps

Theorem (based on [And83], Theorem 8)

Suppose that the value of the primal program is finite. If b is in the interior of $\{Ax \in Z \mid x \geq 0\}$ and on the pre-image of some neighborhood of b in X the value function $\langle x, c \rangle_{XY}$ is bounded then there is no duality gap for (EP).

Brief Overview

- [BHP13] Beiglböck et al. (2013) *Model-Independent Bounds for Option Prices: a Mass Transport Approach*;
- [GHT14] Galichon et al. (2014) *Stochastic Control Approach to No-Arbitrage Bounds Given Marginals*;
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 - Assumption is that marginals are uniquely determined and calibrated to the market;
 - Relaxing the assumption of full marginals: Davis et al. (2013) *Arbitrage Bounds for Prices of Weighted Variance Swaps* [DOR13].

Problem set-up

- Assumptions on the market:
 - ▷ Market is frictionless;
 - ▷ No interest rates, no dividends;
 - ▷ Two time periods t_1 and t_2 ;
 - ▷ Allowed to trade dynamically in the underlying;
 - ▷ Buy and hold positions in other hedging instruments;
- European call options maturing at t_1 $\{k_{1,i}, p_{1,i}\}_{i=1}^{n_1}$ and t_2 $\{k_{2,i}, p_{2,i}\}_{i=1}^{n_2}$ with $n_1, n_2 < \infty$ satisfying no-arbitrage conditions [DH07, Theorem 4.2, p.9]

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Problem set-up

- $\Phi : \mathbb{R}_+^2 \rightarrow \mathbb{R}_+$ is a Borel measurable function that denotes the pay-off of an exotic option at t_2 ;
- $\delta : \mathbb{R}_+ \rightarrow \mathbb{R}$ is a continuous and bounded function in $C_b(\mathbb{R}_+)$ that denotes the delta hedge at time t_1 ;
- $c = (1, 1, p_{1,1}, \dots, p_{1,n_1}, p_{2,1}, \dots, p_{2,n_2})^T \in \mathbb{R}^m$ is a vector of today's prices with $m := 2 + n_1 + n_2$;
- $a(x_1, x_2) = (1, x_2, (x_1 - k_{1,1})_+, \dots, (x_1 - k_{1,m_1})_+, \dots, (x_2 - k_{2,n_2})_+)$ for all $(x_1, x_2) \in \mathbb{R}_+^2$ is a vector of pay-offs.

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Primal Problem (sub-hedging)

$$\underline{\mathcal{P}} := \sup_{\pi \in \underline{\Pi}} \langle c, \pi \rangle,$$

where

$$\underline{\Pi} := \left\{ \pi \in P_+^* \mid \exists \delta \in C_b(\mathbb{R}_+) \text{ s.t. } \Theta^{\pi, \delta}(x_1, x_2) \leq \Phi(x_1, x_2) \right\},$$

the inequality holds for $(x_1, x_2) \in \mathbb{R}_+^2$.

$$\Theta^{\pi, \delta}(x_1, x_2) := A(x_1, x_2)\pi + \delta(x_1)(x_2 - x_1).$$

The linear map A is defined by

$$A(x_1, x_2)\pi := a(x_1, x_2)^T \pi.$$

Primal Problem (super-hedging)

Equally can define a super-hedging problem

$$\bar{\mathcal{P}} := \inf_{\pi \in \bar{\Pi}} \langle c, \pi \rangle,$$

where

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the inequality holds for $(x_1, x_2) \in \mathbb{R}_+^2$.

Primal Problem

Decision variable π in the primal problem is in P_+^*

$$P_+^* = \{\pi \in \mathbb{R}^m \mid \langle x, \pi \rangle \geq 0, \text{ for all } x \in P_+\},$$

where P_+ is

$$P_+ = \{x \in \mathbb{R}^m \mid x = \lambda c, \text{ where } c \in P_+ \text{ for all } \lambda \in \mathbb{R}_+\}.$$

Lemma

The dual cone $P_+^ \in \mathbb{R}^m$ is closed in the usual topology on \mathbb{R}^m if the prices $\{p_{1,i}\}_{i=1}^{n_1}$ and $\{p_{2,i}\}_{i=1}^{n_2}$ of European call options are consistent with absence of arbitrage.*

Remark ([And83])

When the cone of primal decision variables P_+ is closed the primal-dual program system becomes symmetric, i.e. the dual of the dual program is itself an LP and is equal to the primal problem.

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Dual Problem

Sub-hedging dual problem:

$$\underline{\mathcal{D}} := \inf_{\mathbb{Q} \in \mathcal{M}} \int_{\mathbb{R}_+^2} \Phi(x_1, x_2) \mathbb{Q}(dx_1, dx_2),$$

super-hedging dual problem:

$$\overline{\mathcal{D}} := \sup_{\mathbb{Q} \in \mathcal{M}} \int_{\mathbb{R}_+^2} \Phi(x_1, x_2) \mathbb{Q}(dx_1, dx_2),$$

where

$$\mathcal{M} := \left\{ \mathbb{Q} \in \mathcal{Q}_+ \mid A^* \mathbb{Q} = c, \int_{\mathbb{R}_+^2} \delta(x_1)(x_2 - x_1) \mathbb{Q}(dx_1, dx_2) = 0 \right\},$$

and \mathcal{Q}_+ denotes the set of all positive finite regular Borel measures on \mathbb{R}_+^2 .

The adjoint map A^* is defined by

$$A^* \mathbb{Q} := \int_{\mathbb{R}_+^2} a(x_1, x_2) \mathbb{Q}(dx_1, dx_2)$$

Martingale Condition

- A measure is a martingale measure if and only if it satisfies the condition

$$\int_{\mathbb{R}_+^2} \mathbf{1}_{\{x_1 \in A\}} (x_2 - x_1) \mathbb{Q}(dx_1, dx_2) = 0,$$

for all Borel sets $A \in \mathcal{B}(\mathbb{R}_+)$.

- It can be extended to all functions $f \in C_b(\mathbb{R}_+)$

$$\int_{\mathbb{R}_+^2} f(x_1)(x_2 - x_1) \mathbb{Q}(dx_1, dx_2) = 0.$$

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Discretization

- Restrict the domain to $K = K_1 \times K_2 \subset \mathbb{R}_+^2$ s.t. K_1 and K_2 are compact;
- Let \mathcal{S}_n be the set of all available delta hedge strategies such that $\mathcal{S}_n := \{f_1, \dots, f_n\} \subset C(K_1)$ for $n \in \mathbb{N}$;
- $\delta \in \text{Span}(\mathcal{S}_n)$ is a possible delta hedge strategy such that $\delta := \sum_{i=1}^n \lambda_i f_i$ for some $\lambda_i \in \mathbb{R}$ for all $i = 1, \dots, n$.

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Discretized problems

Re-formulate the problem such that for each $n \in \mathbb{N}$ the primal programs read

$$\underline{P}_n := \sup_{\pi \in \underline{\Pi}_n} \langle c, \pi \rangle, \quad (1)$$

where

$$\underline{\Pi}_n := \left\{ \pi \in P_+^* \mid \exists \delta \in \text{Span}(\mathcal{S}_n) \text{ s.t. } \Theta^{\pi, \delta}(x_1, x_2) \leq \Phi(x_1, x_2) \right\}.$$

$$\overline{P}_n := \inf_{\pi \in \overline{\Pi}_n} \langle c, \pi \rangle, \quad (2)$$

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Discretized problems

The dual programs then read

$$\underline{D}_n := \inf_{\mathbb{Q} \in \mathcal{M}_n} \int_K \Phi(x_1, x_2) \mathbb{Q}(dx_1, dx_2), \quad (3)$$

$$\overline{D}_n := \sup_{\mathbb{Q} \in \mathcal{M}_n} \int_K \Phi(x_1, x_2) \mathbb{Q}(dx_1, dx_2), \quad (4)$$

where

$$\mathcal{M}_n := \left\{ \mathbb{Q} \in \mathcal{Q}_+ \mid A^* \mathbb{Q} = c, \int_K f(x_1)(x_2 - x_1) \mathbb{Q}(dx_1, dx_2) = 0 \right\},$$

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for all $f \in \mathcal{S}_n$

Strong duality - application of interior point condition

Theorem (Strong duality for sub-hedging problem)

Let the pay-off function $\Phi : K \rightarrow \mathbb{R}$ be a lower semi-continuous function and assume there exists a constant $C > 0$ such that

$$\Phi(x_1, x_2) \geq -C(1 + |x_1| + |x_2|), \quad \text{for all } x_1, x_2 \geq 0.$$

Assume that the value of dual program is finite and c lies in the interior of the set

$$V_{m+1}^n := \left\{ b \in \mathbb{R}^{m+1} \mid \int_K (a(x_1, x_2), f(x_1)(x_2 - x_1)) \mathbb{Q}(dx_1, dx_2) = b \right\},$$

where $\mathbb{Q} \in \mathcal{Q}_+$ and for all $f \in S_n$. Then the strong duality holds for each $n \in \mathbb{N}$, i.e. $\underline{\mathcal{P}}_n = \underline{\mathcal{D}}_n$. Moreover there exists an optimal portfolio π_* .

Strong duality - application of interior point condition

Corollary (Strong duality for super-hedging problem)

Let the pay-off function $\Phi : K \rightarrow \mathbb{R}$ be an upper semi-continuous function and assume there exists a constant $C > 0$ such that

$$\Phi(x_1, x_2) \leq C(1 + |x_1| + |x_2|), \quad \text{for all } x_1, x_2 \geq 0.$$

*Then the strong duality holds for each $n \in \mathbb{N}$, i.e. $\overline{\mathcal{P}}_n = \overline{\mathcal{D}}_n$.
Moreover there exists an optimal portfolio π^* .*

Limiting case

Lemma

As n tends to infinity, the following limits exist

$$\mathcal{M} = \bigcup_{n=1}^{\infty} \bigcap_{k=n}^{\infty} \mathcal{M}_k = \bigcap_{n=1}^{\infty} \bigcup_{k=n}^{\infty} \mathcal{M}_k,$$

$$\underline{\pi} = \bigcup_{n=1}^{\infty} \bigcap_{k=n}^{\infty} \underline{\pi}_k = \bigcap_{n=1}^{\infty} \bigcup_{k=n}^{\infty} \underline{\pi}_k.$$

$$\bar{\pi} = \bigcup_{n=1}^{\infty} \bigcap_{k=n}^{\infty} \bar{\pi}_k = \bigcap_{n=1}^{\infty} \bigcup_{k=n}^{\infty} \bar{\pi}_k.$$

Attainment of optimal solutions

Lemma

Denoting the space of Borel probability measures on the set K as $P(K)$, \mathcal{M} is closed subset of $P(K)$. Moreover it is compact.

Proposition

If the value of the dual problem \underline{D} (\overline{D} resp.) is finite, then the optimal value is attained and there exists an optimal measure Q_ (Q^* resp.).*

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Further work

- Relaxing the assumption on compactness (i.e. setting $K = \mathbb{R}_+^2$);
- Extending the framework to multiple time periods and then to continuous time;
- Uniqueness of optimal solutions;
- Describing optimal portfolio weights explicitly.

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THANK YOU